# Numerical analysis (1/7): Errors in numerical analysis University of Luxembourg

### Philippe Marchner

Siemens Digital Industries Software, France

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- 1. Introduction: types of errors
- 2. Computer representation of real numbers
- 3. Floating-point arithmetic
- 4. Conditioning and stability
- 5. Summary

# Outline

## 1. Introduction: types of errors

2. Computer representation of real numbers

- 3. Floating-point arithmetic
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What type of errors do we find in numerical analysis ?

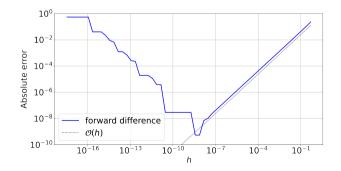
- 1. Errors in the problem to be solved
  - errors in the mathematical model
  - errors in the input data: noise, measurement
- 2. Round-off errors (today's topic, computer floating-point representation)
- 3. Approximation errors
  - Discretization/truncation errors: going from the continuous to the discrete level
  - Convergence errors: iterative methods

## Round-off vs discretization error: example

Derivative approximation (see exercise)

$$f'(x_0) = \frac{f(x_0+h) - f(x_0)}{h} + \mathcal{O}(h)$$

The **discretization error** is linear with hBut when h is too small, we have round-off errors





# 2. Computer representation of real numbers

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A computer has a finite capacity and cannot store all real numbers ( $\pi \approx 3.14159...$ ) Any number is approximated by a rational number, and has a finite number of digits  $d_i$ 

#### Floating point representation of $x \in \mathbb{R}$

$$\mathsf{fl}(x) = \pm m imes b^e, \quad m = \left( rac{d_0}{b^0} + rac{d_1}{b^1} + \dots + rac{d_{p-1}}{b^{p-1}} 
ight)$$

m: mantissa or significand, e: exponent, b: radix or basis, p: precision Normalized (unique) representation:  $1 \le m < b$ ,  $d_0 \ne 0$ 

#### Examples

Let us choose fl(x) = +152853.50, b = 10, p = 8, e = 5152853.50 =  $1.5285350 \times 10^5 = \frac{15285350}{10^7} \times 10^5 = (1 \times 10^0 + 5 \times 10^{-1} + \dots + 0 \times 10^{-7}) \times 10^5$ 

# Floating-point representation

#### Examples

Approximate  $\pi$  in base 2, p = 24:  $\pi \approx 110010010000111111011011$ Approximate 0.1 in base 2, p = 16:  $0.1 \approx 1100110011001101$ 

 $\mathsf{fl}(\pi) = \left(1 + 1 \times 2^{-1} + \dots + 1 \times 2^{-23}\right) \times 2^1 = 3.141592$ **7**, 7 digits precision (why ?)

### Floating point system

A floating point system is characterized by 4 values: (b, p, U, L) base, precision, exponent upper and lower bound such as  $L \le e \le U$ 

Double precision format (64 bit storage)

b = 2, p = 52, L = -1022, U = 1023

11 bit exponent, 52+1 bits for the significand (1 is implicit for the sign)

Let us look what Python uses...

#### Example

Single precision: p = 23 bits significand, 8 bit exponent, U = 127

The sequence of digits

### $1 \ 0111 \ 1110 \quad 100 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000$

considered as simple precision can be interpreted as follows:

- the first digit is 1, so the sign is negative
- the next 8 digits form the exponent: (0111110)<sub>2</sub> =  $0 \times 2^7 + 2^6 + 2^5 + 2^4 + 2^3 + 2^2 + 2^1 + 0 \times 2^0 = 126$
- the next 23 digits form the significand, where the first bit is 1 by convention

The number in decimal is:

 $-2^{(126-U)} \times (1 + 1 \times 2^{-1} + 0 \times 2^{-2} + \dots + 0 \times 2^{-23}) = -0.5 \times 1.5 = -0.75$ 

When rounding, two adjacent numbers are spaced by  $\eta = b^{-p}/2$ , it defines the machine precision.

### Round-off error

Absolute round-off error:  $|fl(x) - x| \le \eta b^e$ Relative round-off error:  $\left|\frac{fl(x)-x}{x}\right| \le \eta$ 

The spacing of a floating point system is constant in the relative sense

### Overflow-underflow

If the exponent e is such as e > U, we have overflow, and if e < L we have underflow

Question: what are the minimum and maximum representable numbers in double precision ?



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# Operations in floating point arithmetic

The usual calculus rules are altered with the floating point representation

Floating operation

$$x \bullet y = (x \bullet y)(1 + \varepsilon), \quad \bullet = (+, -, \times, \div), \quad |\varepsilon| \le \eta$$

Addition is not associative in this context !  $(a + b) + c \neq a + (b + c)$ The relative errors are  $\frac{a+b}{a+b+c}\varepsilon_1(1+\varepsilon_2) + \varepsilon_2$  and  $\frac{b+c}{a+b+c}\varepsilon_3(1+\varepsilon_4) + \varepsilon_4$ 

#### Examples

Derive the relative errors of the above additions For p = 8 try  $a = 2.3371258 \times 10^{-5}$ ,  $b = 3.3678429 \times 10^{1}$ ,  $c = -3.3677811 \times 10^{1}$ 

#### Examples

Compute  $Q = \frac{\pi - 3.1415}{10^4(\pi - 3.1515) - 0.927}$  by truncating  $\pi$  with p significant digits



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Stability is an important concept in numerical analysis, and is found in different contexts:

- stability of numerical schemes (e.g. ODEs, finite differences),
- stability of numerical algorithms (e.g. linear systems),
- stability has also meanings at the continuous level (PDE, dynamical systems),

It is linked to the propagation of errors during the computation.

# Stability

### Forward and backward errors

- The forward error measures the difference between the computed and exact value  $|\hat{y}-y|$
- The backward error is the error on the input:  $|\hat{x} x|$ , for a perturbed output  $\hat{y} = f(\hat{x})$

If the backward error is "small", we say the algorithm to be backward-stable

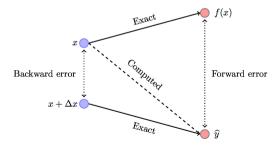


Figure: Illustration from Nicholas J. Higham blog

# Conditioning

The condition number is the ratio of a relative change in the output to a relative change in the input

#### Condition number $\kappa$

$$\kappa = \frac{|(f(\hat{x}) - f(x))/f(x)|}{|(\hat{x} - x)/x|} = \frac{|(\hat{y} - y)/y|}{|(\hat{x} - x)/x|} = \frac{|\Delta y/y|}{|\Delta x/x|}$$

We have the relation

 $|relative forward error| = Condition number \times |relative backward error|$ 

It represents the sensitivity related to the input data.

Conditioning depends on the problem, stability depends on the algorithm

### Condition number of a $C^1$ function f

when  $\hat{x} \approx x$ , we have  $\kappa_f \approx \left| \frac{xf'(x)}{f(x)} \right|$ 

#### Examples

Find the conditioning of the functions  $\sqrt{x}$ , a - x, tan(x)

### Examples

Evaluate the stability of  $f(x) = \sqrt{x+1} - \sqrt{x}$  by computing f(12345) with p = 6 significant digits Try again using the formula  $f(x) = \frac{1}{\sqrt{x+1}+\sqrt{x}}$ Compare the stability properties of the two formulas



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- All machines use a floating-point representation,
- We must be very careful during computations to avoid round-off errors,
- Conditioning and stability analysis are useful theoretical tools,
- Designing numerically stable algorithms is in general not trivial